Goldon-Thompson via pinching inequality

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Given A, B hermitian matrices, Golden-Thompson inequality $[1, 2]$ $[1, 2]$ $[1, 2]$ states that:

$$
tr [exp ((A + B))] \leq tr [exp ((A)) exp ((B))]
$$

It is trivial if A, B commute i.e. $AB = BA$. A quick example would be $A = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ $0 -1$) and $B = \mathbb{I}_2$, while it does not hold for $A = \begin{pmatrix} 1 & 1 \\ 0 & 2 \end{pmatrix}$ and $B = \begin{pmatrix} 1 & 2 \\ 2 & 3 \end{pmatrix}$ as $AB \neq BA$.

There have been multiple proofs $[3]$, but they are not so straight forward to follow. Sutter *et al.* $[4]$ presented a more intutive proof using spectral pinching. This document summarizes their approach and at places is slightly more elaborated than the original version that appeared in [\[4\]](#page-3-3). We will stick to real matrices here.

Spectral Pinching Method

Consider a square complex matrix A partitioned as a $r \times r$ block matrix: $A =$ $\sqrt{ }$ $\overline{}$ A_{11} A_{12} \cdots A_{1r} A_{21} A_{22} \cdots A_{rr} A_{r1} A_{r2} \cdots A_{rr} \setminus $\Bigg\}$

We can decompose this matrix intro two matrices comprising the diagonal and the off-diagonal elements respectively.

$$
A = A_D + A_{\tilde{D}}
$$

\n
$$
A_D = \begin{pmatrix} A_{11} & & & \\ & A_{22} & & \\ & & \ddots & \\ & & & A_{rr} \end{pmatrix}
$$

\n
$$
A_{\tilde{D}} = \begin{pmatrix} 0 & A_{12} & \cdots & A_{1r} \\ A_{21} & 0 & \cdots & A_{2r} \\ \vdots & \ddots & \ddots & \vdots \\ A_{r1} & A_{r2} & \cdots & 0 \end{pmatrix}
$$

 A_D is called a *pinching* of A. The simplest case is that of A_{ij} being 1 dimensional that we will be using here.

Any positive semi-definite matrix A can be written as $A = \sum_{i=1}^{n} \lambda_i P_{\lambda_i}$ where λ_i are n distinct eigen values of A. P_{λ_i} are orthogonal projectors such that $\sum_{i=1}^n P_{\lambda_i} = \overline{\mathbb{I}}$ and hence $P_{\lambda_i}^2 = P_{\lambda_i}$

The *spectral pinching map* of A is then given by:

$$
\mathcal{P}_A: X \mapsto \sum_{\lambda} P_{\lambda} X P_{\lambda}
$$

The entire idea here is to use some form of convex combination resulting in an averaging operation. The pinching map in turn has the following properties:

- (i) $\mathcal{P}_A[X]A = A\mathcal{P}_A[X]$
- (ii) $tr[\mathcal{P}_A[X]A] = tr[XA]$
- (iii) $\mathcal{P}_A[X] \geq \frac{1}{n}X$

Lemma 1. $\mathcal{P}_A[X]A = A\mathcal{P}_A[X]$

Proof:

$$
P_{\lambda_i}^2 = P_{\lambda_i}
$$
 and $P_{\lambda_i} \perp P_{\lambda_j}$

$$
\mathcal{P}_A[X]A = \sum_{i=1}^n P_{\lambda_i} X P_{\lambda_i} \sum_{i=1}^n \lambda_i P_{\lambda_i}
$$

\n
$$
= \sum_{i=1}^n \sum_{j=1}^n \lambda_j P_{\lambda_i} X P_{\lambda_i} P_{\lambda_j}
$$

\n
$$
= \sum_{i=1}^n \lambda_i P_{\lambda_i} X P_{\lambda_i} P_{\lambda_i}
$$

\n
$$
= \sum_{i=1}^n \lambda_i P_{\lambda_i} X P_{\lambda_i}
$$

\n
$$
= \sum_{i=1}^n \lambda_i P_{\lambda_i} P_{\lambda_i} X P_{\lambda_i}
$$

\n
$$
= \sum_{i=1}^n \sum_{j=1}^n \lambda_j P_{\lambda_j} P_{\lambda_i} X P_{\lambda_i}
$$

\n
$$
= \sum_{j=1}^n \lambda_j P_{\lambda_j} \sum_{i=1}^n P_{\lambda_i} X P_{\lambda_i}
$$

\n
$$
= A \mathcal{P}_A[X]
$$

Lemma 2. $tr[\mathcal{P}_A[X]A] = tr[XA]$ Proof:

$$
\operatorname{tr}[\mathcal{P}_A[X]A] = \operatorname{tr}[\sum_{i=1}^n P_{\lambda_i} X P_{\lambda_i} \sum_{i=1}^n \lambda_i P_{\lambda_i}]
$$

$$
= \operatorname{tr}[\sum_{i=1}^n \lambda_i P_{\lambda_i} X P_{\lambda_i}]
$$

$$
\operatorname{tr}[P_{\lambda_i} X P_{\lambda_i}] = \sum_j (P_{\lambda_i} X)_{jj}
$$

$$
\implies \operatorname{tr}[\mathcal{P}_A[X]A] = \sum_{i=1}^n \lambda_i \sum_j (P_{\lambda_i} X)_{jj}
$$

$$
= \operatorname{tr}[AX]
$$

Lemma 3. $\mathcal{P}_A[X] \geq \frac{1}{n}X$

Proof:

Proving this part is probably the trickiest among the four lemmata here, but is the entire key behind deducing the final Golden-Thompson inequality. Consider a unitary matrix U_y defined as $U_y = \sum_{u=1}^{n} e^{i2\pi y u/n} P_{\lambda_u}$. It is easy to verify that $U_y U_y^* = \mathbb{I}$ as $\sum_{i=1}^n P_{\lambda_i} = \mathbb{I}$. Also $U_y \ge 0$ and $U_n = \mathbb{I}$

$$
\sum_{y=1}^{n} U_y X U_y^* = \sum_{y=1}^{n} \sum_{s=1}^{n} \sum_{t=1}^{n} e^{i2\pi y s/n} P_{\lambda_s} X P_{\lambda_t} e^{-i2\pi y t/n}
$$

\n
$$
= \sum_{s=1}^{n} \sum_{t=1}^{n} P_{\lambda_s} X P_{\lambda_t} \sum_{y=1}^{n} e^{i2\pi y (s-t)/n}
$$

\n
$$
= \sum_{s=1}^{n} \sum_{t=1}^{n} P_{\lambda_s} X P_{\lambda_t} (n \mathbf{1}_{\{s=t\}})
$$

\n
$$
= n \sum_{s=1}^{n} P_{\lambda_s} X P_{\lambda_s}
$$

\n
$$
\implies P_A[X] = \sum_{s=1}^{n} P_{\lambda_s} X P_{\lambda_s} = \frac{1}{n} \sum_{y=1}^{n} U_y X U_y^*
$$

\n
$$
\geq \frac{1}{n} X
$$

Now once (iii) is proved, the rest of the steps for proving the GT are straightforward. For a semi-positive definite $d \times d$ matrix A, we have the following lemma.

Lemma 4. $|spec(A^{\bigotimes m})| \leq O(\text{poly}(m))$ Proof:

The number of eigen values for $A^{\bigotimes m}$ is bounded by the number of the number of possible possible combinations of a sequence of d sy, bols (the maximum possible distinct eigen values of A) of length m which is given by $\binom{m+d-1}{d-1} \le \frac{(m+d-1)^{d-1}}{(d-1)!} = O(poly(m))$

Golden Thompson Inequality

Given positive definite matrices A, B and using the facts that $\exp()$ and $\text{tr}[\exp()]$ are operator monotone:

$$
\log tr[\exp(\log A + \log B)] = \frac{1}{m} \log tr[\exp(\log A^{\otimes m} + \log B^{\otimes m})]
$$

\n
$$
\leq \frac{1}{m} \log tr[\exp(\log P_{B^{\otimes m}}[A^{\otimes m}]|\sec(A^{\otimes m})| + \log B^{\otimes m})]
$$
Using Lemma 3
\n
$$
= \frac{1}{m} \log tr[\exp(\log P_{B^{\otimes m}}[A^{\otimes m}] + \log|\sec(A^{\otimes m})| + \log B^{\otimes m})]
$$

\n
$$
= \frac{1}{m} \log tr[\exp(\log P_{B^{\otimes m}}[A^{\otimes m}] + \log B^{\otimes m})] + \frac{1}{m}|\sec(A^{\otimes m})|
$$

\n
$$
\leq \frac{1}{m} \log tr[\exp(\log P_{B^{\otimes m}}[A^{\otimes m}] + \log B^{\otimes m})] + \frac{O(\text{poly}(m))}{m}
$$
Using Lemma 4
\n
$$
= \frac{1}{m} \log tr[\exp(\log P_{B^{\otimes m}}[A^{\otimes m}]B^{\otimes m})] + \frac{O(\text{poly}(m))}{m}
$$
Using Lemma 1
\n
$$
= \frac{1}{m} \log tr[P_{B^{\otimes m}}[A^{\otimes m}]B^{\otimes m}] + \frac{O(\text{poly}(m))}{m}
$$
Using Lemma 1
\n
$$
= \frac{1}{m} \log tr[A^{\otimes m}B^{\otimes m}] + \frac{O(\text{poly}(m))}{m}
$$
Using Lemma 2
\n
$$
= \log tr[AB] + \frac{O(\text{poly}(m))}{m}
$$

$$
\implies \operatorname{tr}[\exp(\log A + \log B)] \le \operatorname{tr}(AB)
$$

$$
\implies \operatorname{tr}[\exp(A + B)] \le \operatorname{tr}(\exp(A) \exp(B))
$$

References

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- [4] D. Sutter, M. Berta, and M. Tomamichel, "Multivariate trace inequalities," Communications in Mathematical Physics, vol. 352, no. 1, pp. 37–58, 2017.